

Binomial arithmetical rank of toric ideals associated to graphs

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Binomial ideals

- Consider the polynomial ring $K[x_1, \dots, x_m]$ over an algebraically closed field K .
- $\mathbf{x}^{\mathbf{u}} := x_1^{u_1} \cdots x_m^{u_m}$ for $\mathbf{u} = (u_1, \dots, u_m) \in \mathbb{N}^m$.
- **Binomial** $B = \mathbf{x}^{\mathbf{u}} - \mathbf{x}^{\mathbf{v}}$ is a difference of two monomials.
- **Binomial ideal** is an ideal generated by binomials.

Definition of a toric ideal of a graph

- G finite, connected and undirected graph having no loops and no multiple edges with vertices $\{v_1, \dots, v_n\}$ and edges $\{e_1, \dots, e_m\}$.
- To $e = \{v_i, v_j\}$ we associate a monomial $\mathbf{t}^e = t_i t_j \in K[t_1, \dots, t_n]$.
- The **toric ideal** I_G is the kernel of the K -algebra homomorphism $\phi : K[x_1, \dots, x_m] \longrightarrow K[t_1, \dots, t_n]$ given by

$$\phi(x_i) = \mathbf{t}^{e_i}, \text{ for every } i = 1, \dots, m.$$

Toric ideals are binomial ideals

- To $e = \{v_i, v_j\}$ we associate a vector $\mathbf{a}_e = (0, \dots, 0, 1, 0, \dots, 0, 1, 0, \dots, 0) \in \mathbb{Z}^n$.
- The toric ideal I_G is generated by all the binomials

$$x_1^{u_1} \cdots x_m^{u_m} - x_1^{v_1} \cdots x_m^{v_m}$$

such that

$$u_1 \mathbf{a}_1 + \cdots + u_m \mathbf{a}_m = v_1 \mathbf{a}_1 + \cdots + v_m \mathbf{a}_m.$$

Example

- G graph with 8 vertices and edges $e_1 = \{1, 2\}$, $e_2 = \{2, 3\}$, $e_3 = \{3, 4\}$, $e_4 = \{4, 5\}$, $e_5 = \{5, 6\}$, $e_6 = \{4, 6\}$, $e_7 = \{4, 7\}$, $e_8 = \{7, 8\}$, $e_9 = \{3, 8\}$, $e_{10} = \{1, 3\}$.
- I_G is minimally generated by

$$x_3x_8 - x_7x_9, x_1x_3^2x_5 - x_2x_4x_6x_{10}.$$

- $\Gamma = (e_{i_1}, \dots, e_{i_{2q}})$ even closed walk of G .

$$f_\Gamma = \prod_{k=1}^q x_{i_{2k-1}} - \prod_{k=1}^q x_{i_{2k}} \in I_G.$$

- I_G is generated by all the binomials f_Γ .

Arithmetical rank

- The **arithmetical rank** $\text{ara}(I_G)$ of I_G is the smallest integer s for which there exists polynomials F_1, \dots, F_s in I_G such that $\sqrt{I_G} = \sqrt{(F_1, \dots, F_s)}$.
- G graph with 5 vertices and edges $e_1 = \{1, 2\}$, $e_2 = \{2, 3\}$, $e_3 = \{3, 4\}$, $e_4 = \{4, 5\}$, $e_5 = \{1, 5\}$, $e_6 = \{1, 3\}$, $e_7 = \{2, 5\}$, $e_8 = \{3, 5\}$.
- I_G is minimally generated by

$$x_1x_8 - x_6x_7, x_2x_5 - x_6x_7, x_3x_5 - x_4x_6, x_2x_4 - x_3x_7.$$

Binomial arithmetical rank

- 1 The **binomial arithmetical rank** $\text{bar}(I_G)$ of I_G is the smallest integer s for which there exists binomials B_1, \dots, B_s in I_G such that

$$\sqrt{I_G} = \sqrt{(B_1, \dots, B_s)}.$$

What is the exact value of $\text{bar}(I_G)$?

- If I_G is complete intersection, then $\text{bar}(I_G) = \text{ht}(I_G)$.
- For the ideal $J = (x_1x_6 - x_3x_4, x_1x_5 - x_2x_4, x_3x_5 - x_2x_6)$ it is known that $\text{ara}(J) = 3$, so $\text{bar}(J) = 3$.
- $J = I_G$ for G the bipartite graph with 5 vertices and edges $e_1 = \{1, 4\}$, $e_2 = \{2, 4\}$, $e_3 = \{3, 4\}$, $e_4 = \{1, 5\}$, $e_5 = \{2, 5\}$, $e_6 = \{3, 5\}$.
- A **bipartite** graph G is one whose vertex set can be partitioned into two subsets X and Y , so that each edge has one vertex in X and one vertex in Y .

Aims of the talk

Theorem

If G is bipartite, then $\text{bar}(I_G) = \mu(I_G)$.

Theorem

If I_G is generated by quadratic binomials, then $\text{bar}(I_G) = \mu(I_G)$.

Circuits of a toric ideal

- The **support of a monomial** $\mathbf{x}^{\mathbf{u}}$ is $\text{supp}(\mathbf{x}^{\mathbf{u}}) = \{i | x_i \text{ divides } \mathbf{x}^{\mathbf{u}}\}$.
- The **support of a binomial** $B = \mathbf{x}^{\mathbf{u}} - \mathbf{x}^{\mathbf{v}}$ is $\text{supp}(B) = \text{supp}(\mathbf{x}^{\mathbf{u}}) \cup \text{supp}(\mathbf{x}^{\mathbf{v}})$.
- An irreducible binomial $B = \mathbf{x}^{\mathbf{u}} - \mathbf{x}^{\mathbf{v}} \in I_G$ is called a **circuit** of I_G if there is no binomial $B' \in I_G$ such that $\text{supp}(B') \subsetneq \text{supp}(B)$.

Example

- G graph with 8 vertices and edges $e_1 = \{1, 2\}$, $e_2 = \{2, 3\}$, $e_3 = \{3, 4\}$, $e_4 = \{4, 5\}$, $e_5 = \{5, 6\}$, $e_6 = \{4, 6\}$, $e_7 = \{4, 7\}$, $e_8 = \{7, 8\}$, $e_9 = \{3, 8\}$, $e_{10} = \{1, 3\}$.
- The circuits are

$$x_3x_8 - x_7x_9, x_1x_3^2x_5 - x_2x_4x_6x_{10},$$

$$x_1x_9^2x_7^2x_5 - x_2x_8^2x_4x_6x_{10}.$$

The graph Δ_G

- $\mathcal{C} = \{E \mid \text{supp}(\mathbf{x}^u) = E \text{ or } \text{supp}(\mathbf{x}^v) = E \text{ where } \mathbf{x}^u - \mathbf{x}^v \text{ is circuit}\}$.
- \mathcal{C}_{\min} is the set of minimal elements of \mathcal{C} .
- The graph Δ_G has vertices the elements of \mathcal{C}_{\min} .
- $\{E, E'\}$ is an edge of Δ_G if and only if there exists a circuit $\mathbf{x}^u - \mathbf{x}^v \in I_G$ such that $\text{supp}(\mathbf{x}^u) = E$ and $\text{supp}(\mathbf{x}^v) = E'$.

Example

- G graph with 5 vertices and edges $e_1 = \{1, 2\}$, $e_2 = \{1, 3\}$, $e_3 = \{1, 4\}$, $e_4 = \{2, 3\}$, $e_5 = \{2, 5\}$, $e_6 = \{3, 4\}$, $e_7 = \{3, 5\}$, $e_8 = \{4, 5\}$.
- The circuits are

$$x_1x_7 - x_2x_5, x_1x_6 - x_3x_4, x_1x_8 - x_3x_5, x_4x_8 - x_5x_6,$$

$$x_3x_7 - x_2x_8, x_1x_6x_7 - x_2x_4x_8, x_2x_5x_6 - x_3x_4x_7.$$

Continue the example

- $\mathcal{C}_{\min} = \{E_1 = \{1, 7\}, E_2 = \{2, 5\}, E_3 = \{1, 6\}, E_4 = \{3, 4\}, E_5 = \{1, 8\}, E_6 = \{3, 5\}, E_7 = \{4, 8\}, E_8 = \{5, 6\}, E_9 = \{3, 7\}, E_{10} = \{2, 8\}\}$.
- The graph Δ_G has 5 connected components, namely

$$\{E_1, E_2\}, \{E_3, E_4\}, \{E_5, E_6\}, \{E_7, E_8\}, \{E_9, E_{10}\}.$$

Definitions

- A subset \mathcal{M} of the edges of a graph G is called a **matching in G** if there are no two edges in \mathcal{M} which are incident with a common vertex.
- \mathcal{M} is a **maximum matching** if it has the maximum possible cardinality among all matchings.
- The cardinality $\tau(G)$ of a maximum matching in G is commonly known as its **matching number**.

Lower bound for $\bar{\nu}$

Definition

For the graph Δ_G we define $\delta(\Delta_G) = \#\text{Vertices}(\Delta_G) - \tau(\Delta_G)$.

Theorem

For a toric ideal I_G we have that

$$\bar{\nu}(I_G) \geq \delta(\Delta_G).$$

Indispensable monomials and binomials

- A binomial $B \in I_G$ is called **indispensable** if every system of binomial generators of I_G contains B or $-B$, while a monomial M is called **indispensable** if every system of binomial generators of I_G contains a binomial B such that the M is a monomial of B .

Bipartite graphs

Theorem

Let G be a bipartite graph, then $\text{bar}(I_G) = \mu(I_G)$.

Sketch of the proof.

- $\mathcal{C}_{\min} = \{\text{supp}(M_1), \dots, \text{supp}(M_r)\}$, where M_1, \dots, M_r are the indispensable monomials.
- $\{E, E'\}$ is an edge of Δ_G if and only if there is an indispensable binomial $M_i - M_j \in I_G$ with $E = \text{supp}(M_i)$ and $E' = \text{supp}(M_j)$.
- Every connected component of Δ_G is an edge.

Complete graph on 4 vertices

- \mathcal{K}_4 complete graph on 4 vertices with edges $e_1 = \{1, 2\}$, $e_2 = \{2, 3\}$, $e_3 = \{3, 4\}$, $e_4 = \{1, 4\}$, $e_5 = \{1, 3\}$, $e_6 = \{2, 4\}$.
- $I_{\mathcal{K}_4}$ is minimally generated by

$$B_1 = x_1x_3 - x_2x_4, B_2 = x_1x_3 - x_5x_6.$$

- The circuits are:

$$B_1, B_2, B_3 = x_2x_4 - x_5x_6.$$

Toric ideals generated by quadratic binomials

Theorem

If I_G is generated by quadratic binomials, then $\text{bar}(I_G) = \mu(I_G)$.

Sketch of the proof.

- $\mathcal{C}_{\min} = \{\text{supp}(M_1), \dots, \text{supp}(M_r)\}$, where M_1, \dots, M_r are the indispensable monomials.
- Δ_G has g connected components which are edges, where g is the number of indispensable binomials in I_G .
- Δ_G has $\frac{t-g}{2}$ connected components which are triangles, where $t = \mu(I_G)$.

Compute $\text{bar}(I_{\mathcal{K}_n})$

- \mathcal{K}_n , $n \geq 4$, is the complete graph on the vertex set $\{v_1, \dots, v_n\}$.
- Every connected component of $\Delta_{\mathcal{K}_n}$ is a triangle.
- $\Delta_{\mathcal{K}_n}$ has $3\binom{n}{4}$ vertices and $\binom{n}{4}$ triangles, corresponding to all complete subgraphs of \mathcal{K}_n of order 4.
- $\text{bar}(I_{\mathcal{K}_n}) = 2\binom{n}{4} = \frac{n(n-1)(n-2)(n-3)}{12}$.
- $\binom{n}{2} - n = \frac{n(n-3)}{2} \leq \text{ara}(I_{\mathcal{K}_n}) \leq \binom{n}{2} = \frac{n(n-1)}{2}$.

G -homogeneous arithmetical rank

- $\deg_G(x_i) = \mathbf{a}_i$ for $i = 1, \dots, m$.
- $\deg_G(\mathbf{x}^{\mathbf{u}}) = u_1 \mathbf{a}_1 + \dots + u_m \mathbf{a}_m$.
- The **G -homogeneous arithmetical rank** $\text{ara}_G(I_G)$ of I_G is the smallest integer s for which there exists G -homogeneous polynomials F_1, \dots, F_s in I_G such that

$$\sqrt{I_G} = \sqrt{(F_1, \dots, F_s)}.$$

- The equality $\text{ara}_G(I_G) = \mu(I_G)$ holds for the two cases studied in this talk.